Exercise 9.9

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i. For each one of the k answers computed, the probability for each one of them being too large is given by $\Pr[\tilde{D} > 2D(\sigma)] = \frac{7}{10}$. If the $\frac{k}{2}$ -smallest answer is reported, then the probability of the answer being too large is

given by the probability of at least $\frac{k}{2}$ answers being too large. This probability is given by $\left(\Pr[\widetilde{D} > 2D(\sigma)]\right)^{\frac{n}{2}} =$

 $\left(\frac{7}{10}\right)^{\overline{2}} > \frac{7}{10}$ for $k \ge 2$.

ii. We should instead report the answer which is in the 'middle' of the 'correct range'. Here, we have that the smallest $\frac{1}{5}$ answers are likely too small, the next $\frac{1}{10}$ answers are likely in the acceptable range, and the highest $\frac{7}{10}$ answers are likely too high.

Thus, we choose the $\frac{1}{5} + \frac{1}{10} = \frac{1}{5} + \frac{1}{20} = \frac{5}{20} = \frac{1}{4}$ th lowest answer.

- iii. See below.
- iv. You would always report the smallest answer; since no answer is too small, the smallest answer is always most likely to fall in the acceptable range.
- v. The new algorithm would give a wrong answer only if all of the *k* answers are wrong. Since the probability of one of the k answers being wrong is equal to $\frac{9}{10}$, the probability of k of them being wrong is $\left(\frac{9}{10}\right)^k$. This means that the success probability is given by $1 - \left(\frac{9}{10}\right)^k$.
- iii. Let X_i denote an indicator random variable which is 1 if the i^{th} answer is too large, and let Y_i denote an indicator random variable which is 1 if the i^{th} answer is too small. Let $X = \sum_{i=1}^{k} X_i$ and $Y = \sum_{i=1}^{k} Y_i$. Now, we know that $E[X_i] = \frac{7}{10}$ and $E[Y_i] = \frac{1}{5}$. By linearity of expectation, we then have that $E[X] = \frac{7k}{10}$ and $E[Y] = \frac{k}{5}$. We then have that the final answer is not a good estimate if either the number of too-large-answers is at least $\frac{3k}{4}$ or if the number of too-small-answers is at least $\frac{k}{4}$. Bounds on these probabilities are given by

$$\Pr\left[X > \frac{3k}{4}\right] = \Pr\left[X > \left(1 + \frac{1}{14}\right) \cdot E[X]\right] < \left(\frac{e^{\frac{1}{14}}}{\left(1 + \frac{1}{14}\right)^{1 + \frac{1}{14}}}\right)^{\frac{7k}{10}}$$
$$\Pr\left[Y > \frac{k}{4}\right] = \Pr\left[X > \left(1 + \frac{1}{4}\right) \cdot E[Y]\right] < \left(\frac{e^{\frac{1}{4}}}{\left(1 + \frac{1}{4}\right)^{1 + \frac{1}{4}}}\right)^{\frac{k}{5}}$$

Then, it follows that the probability of an answer not being a good estimate is bounded by

$$\Pr[answer \ is \ a \ good \ estimate] > 1 - \left(\frac{e^{\frac{1}{14}}}{\left(1 + \frac{1}{14}\right)^{1 + \frac{1}{14}}}\right)^{\frac{7\pi}{10}} - \left(\frac{e^{\frac{1}{4}}}{\left(1 + \frac{1}{4}\right)^{1 + \frac{1}{4}}}\right)^{\frac{\pi}{5}}$$

$$(1+7)\frac{1}{5} = \frac{1}{4}$$
$$\frac{4}{20} + \frac{1\cdot4}{20} = \frac{5}{20}$$

$$? \cdot 4 = 1$$

? = $\frac{1}{4}$

 $\begin{array}{ccc} 1 \\ 1 \\ - \end{array} \end{array}$

(i) answer is not too large \Leftrightarrow at leasthalf of the answers not too large We want to bound Ir [Inumber of answers not too large) $\geq \frac{k}{2}$] define indicator random variables Xi = (if it as wernot too lorge X = Ei= Xi = Haywouthet orienstoo large => compute E[X] and use Chernoff bound $E[X] = E[z_{i}, X_i]$ $= \sum_{i=1}^{\ell} [E G_{Xi}]$ $= \mathcal{E}_{i=1}^{l} \operatorname{Pr}[X_{i=1}]$ $(1+?) \frac{4}{10} = \frac{3}{4}$ $-\frac{1}{2} - \frac{3}{10}$ $\frac{2d}{40} + \frac{2d}{40} =$ $=\frac{3k}{10}$ cherroft bound: $\Pr[X > (i+\delta) \in [X]] < \left(\frac{e^{\delta}}{(i+\delta)^{H_{\delta}}}\right)^{H_{\delta}} = \sum_{i=1}^{2} \frac{3k}{i^{\circ}}$ $\Pr[X > \frac{k}{2}] = \Pr[X > \frac{5}{3} \in [X]] < \left(\frac{e^{\frac{2}{3}}}{(\frac{5}{3})^{\frac{5}{3}}}\right)^{\frac{3k}{i^{\circ}}} < \sum_{i=1}^{2} \frac{1}{i^{\circ}}$ $< (0, dy)^{3 \pm 1/0}$ $\gamma \cdot {}^2 P = 2$ In [answer too large] = 1- In [answer not too large]

 $> 1 - (0 \cdot \partial_{4}) \frac{3 \cdot k}{10}$

the longer h, the greater the probability

> 2un overall i 2 th